





CHAPTER 4: ROUTING DYNAMIC LANJUTAN

Routing & Switching





- 4.1 Dynamic Routing Protocols
- 4.2 Distance Vector Dynamic Routing
- 4.3 RIP and RIPng Routing
- 4.4 Link-State Dynamic Routing
- 4.5 The Routing Table
- 4.6 Summary





Routing Protocols Classification





DYNAMIC ROUTING PROTOCOL OPERATION THE EVOLUTION OF DYNAMIC ROUTING PROTOCOLS

- Dynamic routing protocols used in networks since the late 1980s
- Newer versions support the communication based on IPv6

Routing Protocols Classification

	Interior Gate	way Protocol	Exterior Gateway Protocols		
	Distance Vector		Link-State		Path Vector
IPv4	RIPv2	EIGRP	OSPFv2	IS-IS	BGP-4
IPv6	RIPng	EIGRP for IPv6	OSPFv3	IS-IS for IPv6	BGP-MP



TYPES OF ROUTING PROTOCOLS DISTANCE VECTOR OR LINK-STATE ROUTING PROTOCOLS

Distance vector protocols use routers as sign posts along the path to the final destination.

A link-state routing protocol is like having a complete map of the network topology. The sign posts along the way from source to destination are not necessary, because all link-state routers are using an identical map of the network. A link-state router uses the link-state information to create a topology map and to select the best path to all destination networks in the topology.



TYPES OF ROUTING PROTOCOLS

Link-State Protocol Operation



Link-state protocols forward updates when the state of a link changes.



TYPES OF ROUTING PROTOCOLS CLASSFUL ROUTING PROTOCOLS

Classful routing protocols do not send subnet mask information in their routing updates:

- Only RIPv1 and IGRP are classful.
- Created when network addresses were allocated based on classes (class A, B, or C).
- Cannot provide variable length subnet masks (VLSMs) and classless interdomain routing (CIDR).
- Create problems in discontiguous networks.



TYPES OF ROUTING PROTOCOLS CLASSLESS ROUTING PROTOCOLS

Classless routing protocols include subnet mask information in the routing updates:

- RIPv2, EIGRP, OSPF, and IS_IS
- Support VLSM and CIDR
- IPv6 routing protocols



TYPES OF ROUTING PROTOCOLS ROUTING PROTOCOL CHARACTERISTICS

	Distance	Vector	Link State			
	RIPv1	RIPv2	IGRP	EIGRP	OSPF	IS-IS
Speed Convergence	Slow	Slow	Slow	Fast	Fast	Fast
Scalability - Size of Network	Small	Small	Small	Large	Large	Large
Use of VLSM	No	Yes	No	Yes	Yes	Yes
Resource Usage	Low	Low	Low	Medium	High	High
Implemenation and Maintenance	Simple	Simple	Simple	Complex	Complex	Complex



TYPES OF ROUTING PROTOCOLS ROUTING PROTOCOL METRICS

A metric is a measurable value that is assigned by the routing protocol to different routes based on the usefulness of that route:

- Used to determine the overall "cost" of a path from source to destination.
- Routing protocols determine the best path based on the route with the lowest cost.



TABEL NILAI DEFAULT ADMINISTRATIVE DISTANCE (AD) PADA ROUTER CISCO

Route Source	Default Distance Values		
Connected interface	0		
Static route	1		
Enhanced Interior Gateway Routing Protocol (EIGRP) summary route	5		
External Border Gateway Protocol (BGP)	20		
Internal EIGRP	90		
IGRP	100		
OSPF	110		
Intermediate System-to-Intermediate System (IS-IS)	115		
Routing Information Protocol (RIP)	120		
Exterior Gateway Protocol (EGP)	140		
On Demand Routing (ODR)	160		
External EIGRP	170		
Internal BGP	200		
Unknown*	255		



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A Link-State Routing Algorithm

Dijkstra's algorithm

- net topology, link costs known to all nodes
 - accomplished via "link state broadcast"
 - all nodes have same info
- computes least cost paths from one node ('source'') to all other nodes
 - gives forwarding table for that node
- iterative: after k iterations, know least cost path to k dest.'s

Notation:

- C(x,y): link cost from node
 x to y; = ∞ if not direct
 neighbors
- D(v): current value of cost of path from source to dest. v
- p(v): predecessor node along path from source to v
- N': set of nodes whose least cost path definitively known





DIJKSTRA'S ALGORITHM: EXAMPLE







Resulting shortest-path tree from u:



 $\begin{array}{c|c} \hline \textbf{Resulting forwarding table in u:} \\ \hline \textbf{destination} & \textbf{link} \\ \hline \textbf{v} & (\textbf{u}, \textbf{v}) \\ \textbf{x} & (\textbf{u}, \textbf{x}) \\ \textbf{x} & (\textbf{u}, \textbf{x}) \\ \textbf{y} & (\textbf{u}, \textbf{x}) \\ \textbf{w} & (\textbf{u}, \textbf{x}) \\ \textbf{z} & (\textbf{u}, \textbf{x}) \end{array}$



LINK-STATE ROUTING PROTOCOL OPERATION DIJKSTRA'S ALGORITHM

Dijkstra's Shortest Path First Algorithm

Shortest Path for host on R2 LAN to reach host on R3 LAN: R2 to R1 (20) + R1 to R3 (5) + R3 to LAN (2) = 27





LINK-STATE UPDATES

Link-State Routing Process

- Each router learns about each of its own directly connected networks.
- Each router is responsible for "saying hello" to its neighbors on directly connected networks.
- Each router builds a Link State Packet (LSP) containing the state of each directly connected link.
- Each router floods the LSP to all neighbors who then store all LSP's received in a database.
- Each router uses the database to construct a complete map of the topology and computers the best path to each destination networks.



The first step in the link-state routing process is that each router learns about its own links and its own directly connected networks.

Link-State of Interface Fa0/0



Link 1

- Network: 10.1.0.0/16
- IP address: 10.1.0.1
- Type of network: Ethernet
- Cost of that link: 2
- Neighbors: None

Link-State of Interface S0/0/0



Link 2

- Network: 10.2.0.0/16
- IP address: 10.2.0.1
- Type of network: Serial
- Cost of that link: 20
- Neighbors: R2



WHY USE LINK-STATE ROUTING PROTOCOLS WHY USE LINK-STATE PROTOCOLS?

Disadvantages of Link-State Routing Protocols

- Maintaining a link-state database and SPF tree requires additional memory.
- Calculating the SPF algorithm also requires additional CPU processing.
- Bandwidth can be adversely affected by link-state packet flooding.





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Thank you very much for your kind attention